Comprehensive formal security analysis of Single Sign-On protocols based on the SAML 2.0 standard



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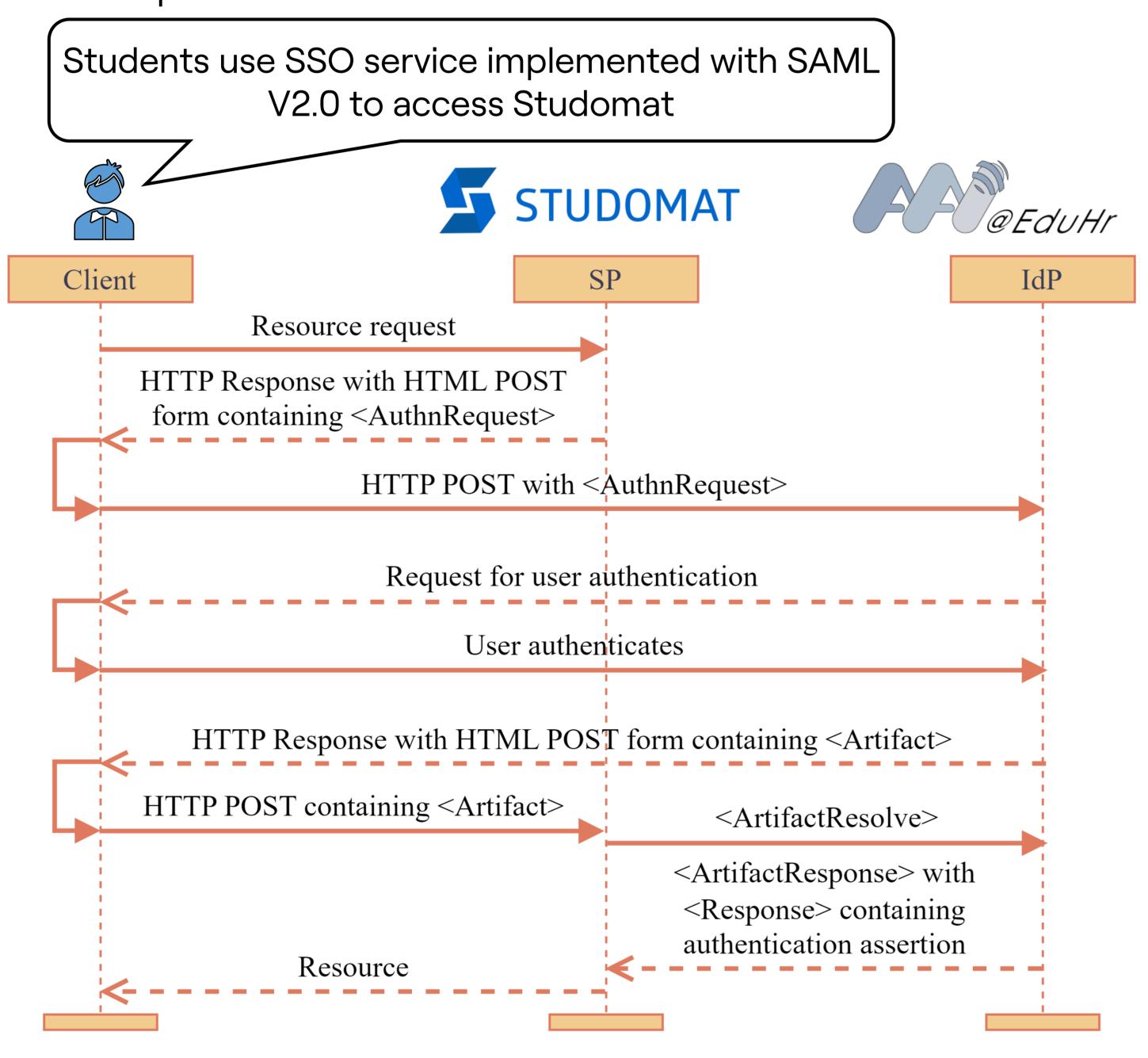
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1. Introduction

(SSO) protocols enable Sign-On Single exchange of authentication and authorization data between trusted service providers and identity providers, enabling user to have one set of credentials at identity provider and being able to use services at multiple service providers. SAML V2.0 [1] is one of the most used protocols that provide SSO functionality. SAML V2.0 is mostly used by governmental, federal and education institutions. Example implementations are Croatian National Identity and Authentication System (NIAS) and Authentication Authorization Infrastructure of Science and Higher Education in Croatia (AAI@EduHr).

2. Problem Description

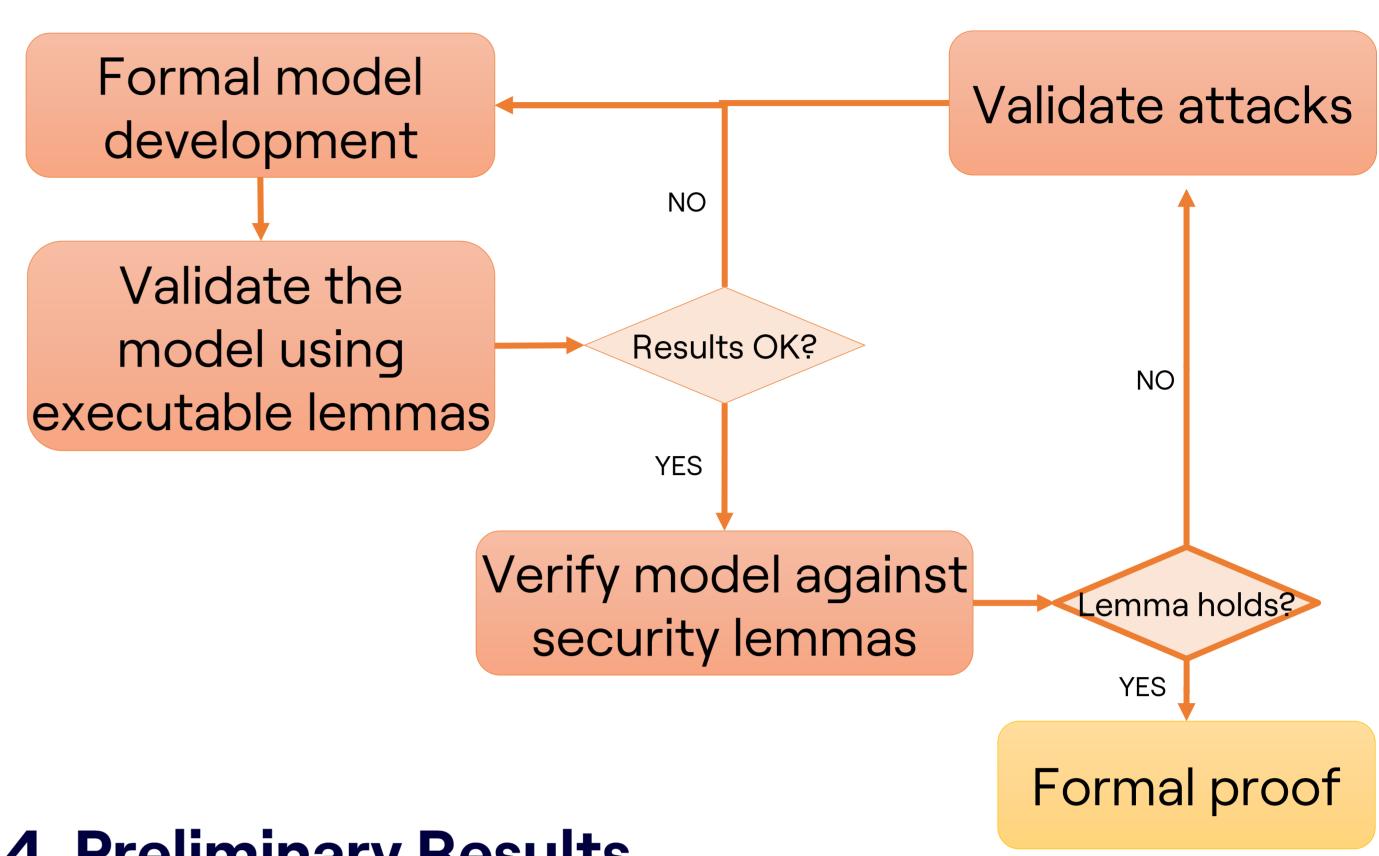
Compromising IdP or braking SSO protocol can allow a malicious actor to access a broad set of services and consequently a broad set of sensitive user data. SSO protocol is single point of failure therefore it would be of much use to have formal proof that it is secure. It is complex to create a formal model of SAML V2.0 SSO Profile since it offers six different use case implementations with various options.



IS THE SAML V2.0 SSO PROTOCOL REALLY SECURE?

3. Methodology

Formal analysis could give us guarantee that SAML V2.0 SSO Profile is secure or can provide us enough knowledge on how to improve the protocol security. **Tamarin Prover** is the state-of-theart tool for formal analysis of security protocols [2]. We used it to create formal model of the use case SP-initiated SSO with POST/Artifact bindings from SAML V2.0 web browser SSO profile. To build such a model we needed to iterate many times through following process.



4. Preliminary Results

For now we managed to formally model one use case SP-initiated SSO with POST/Artifact bindings from SAML V2.0 web browser SSO profile. We generated 8 different variants of use case in accordance with different possible implementations of secure ID generation, RelayState handling and authentication request signing. Results of analysing broad set of security properties of given use case is shown in following table.

No	Property	Protocol variant							
No.		srs	srw	sns	snw	urs	urw	uns	unw
1	sec_Client_Resource_Authenticity	✓	✓	✓	✓	✓	✓	✓	✓
2	sec_Client_Resource_Authenticity_Strong	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
3	sec_Client_Registered_Resource_Secrecy	✓	\checkmark	✓	✓	\checkmark	✓	\checkmark	✓
4	sec_Client_Registered_Resource_Secrecy_Strong	X	X	X	×	X	X	X	X
5	sec_Client_Resource_Secrecy	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	✓
6	sec_Client_Resource_Secrecy_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	✓
7	sec_Client_Resource_Freshness	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	✓
8	sec_Client_SP_Non_Injective_Agreement	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
9	sec_Client_SP_Non_Injective_Agreement_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
10	sec_Client_SP_Injective_Agreement	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
11	sec_Client_SP_Injective_Agreement_Strong	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
12	sec_SP_Client_Resource_Secrecy	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	✓
13	sec_SP_Client_Resource_Secrecy_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	✓
14	sec_SP_Client_Non_Injective_Agreement	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
15	sec_SP_Client_Non_Injective_Agreement_Strong	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
16	sec_SP_IdP_Weak_Agreement	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
17	sec_SP_IdP_Weak_Agreement_Strong	\checkmark	\checkmark	✓	✓	\checkmark	✓	\checkmark	\checkmark
18	sec_SP_IdP_Non_Injective_Agreement	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
19	sec_SP_IdP_Non_Injective_Agreement_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
20	sec_SP_IdP_Authentication_Freshness	\checkmark	\checkmark	\checkmark	✓	\checkmark	X	\checkmark	X
21	sec_SP_IdP_Authentication_Freshness_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	X	\checkmark	X
22	sec_SP_IdP_Assertion_Secrecy	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
23	sec_SP_IdP_Assertion_Secrecy_Strong	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark	\checkmark	\checkmark
24	sec_IdP_SP_Non_Injective_Pre_Agreement	\checkmark	\checkmark	\checkmark	✓	X	X	X	X
25	sec_IdP_SP_Weak_Agreement	\checkmark	\checkmark	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark
26	sec_IdP_SP_Non_Injective_Agreement	X	X	X	X	X	X	X	X
27	sec_IdP_SP_Injective_Agreement	X	X	X	X	X	X	X	X
28	sec_IdP_SP_Assertion_Secrecy	✓	✓	✓	✓	✓	✓	✓	✓

first letter of the protocol variant – authorization requests are signed ("s") or unsigned ("u") second letter of the protocol variant - relay state mechanism is used ("r") or not used ("n") third letter of the protocol variant – random IDs are cryptographically strong ("s") or weak ("w") ✓ – property verified, X – property falsified

5. Conclusion

The results of formal verification of security proprieties of SPinitiated SSO with POST/Artifact bindings have shown that most important security properties of resource and assertion secrecy hold in all possible variants of protocol.

In future work we would like to create comprehensive formal models of all possible use cases of Web Browser SAML SSO Profile and to automatically verify presented security properties.

References

[1] Profiles for the OASIS Security Assertion Markup Language (SAML) V2.0, OASIS Std., March 2005. [Online]. Available: https://docs.oasis-open.org/security/saml/v2.0/saml-profiles-2.0-os.pdf



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